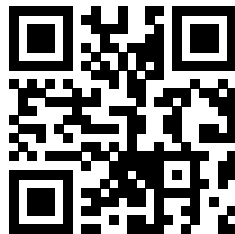


Probabilistic Bijections for Non-Attacking Fillings

Guilherme Zeus Dantas e Moura
(joint work with Olya Mandelshtam)

University of Waterloo

May 24, 2025

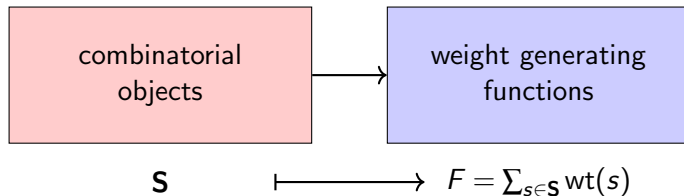


Ontario Combinatorics Workshop 2025

Part I

Probabilistic Bijections

In Combinatorial Enumeration

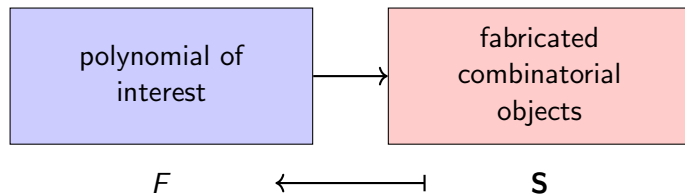


$\text{wt}(s)$ can be $x^{(\text{size of } s)}$, or it can be more general.

Learning about F tells us about $\mathbf{S}$.

- e.g., asymptotics, exact counts, etc.

How Combinatorialists do Algebra?



Learning about S tells us about F .

Simplest goal: $F = G$

F and G are weight generating functions of $\mathbf{S}$ and $\mathbf{T}$. How to combinatorially prove

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F and G are weight generating functions of $\mathbf{S}$ and $\mathbf{T}$. How to combinatorially prove

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A strategy

Construct a **weight-preserving bijection** between $\mathbf{S}$ and $\mathbf{T}$.

Weight-preserving Bijections

Weight-preserving bijections:

$$f: \mathbf{S} \rightarrow \mathbf{T} \quad \text{with an inverse} \quad g: \mathbf{T} \rightarrow \mathbf{S}$$

such that, whenever $f(S) = T$, we have

$$\text{wt}(S) = \text{wt}(T).$$

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Proposition

If there exists a weight-preserving bijection between $\mathbf{S}$ and $\mathbf{T}$, then

$$\sum_{S \in \mathbf{S}} \text{wt}(S) = \sum_{T \in \mathbf{T}} \text{wt}(T).$$

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What about when it's hard (or impossible) to find a weight-preserving bijection?

Probabilistic Bijections

Probabilistic bijection: pair of maps to an algebra A (Frieden and Schreier-Aigner 2024)

$$\text{prob}_{\mathbf{S}}: \mathbf{S} \times \mathbf{T} \rightarrow A \quad \text{and} \quad \text{prob}_{\mathbf{T}}: \mathbf{T} \times \mathbf{S} \rightarrow A$$

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Proof of Proposition

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Proof

$$\sum_{S \in \mathbf{S}} 1 \cdot \text{wt}(S)$$



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Proof

$$\sum_{S \in \mathbf{S}} 1 \cdot \text{wt}(S) = \sum_{S \in \mathbf{S}} \sum_{T \in \mathbf{T}} \text{prob}_{\mathbf{S}}(S, T) \text{wt}(S) \quad (\text{probabilities sum to 1})$$



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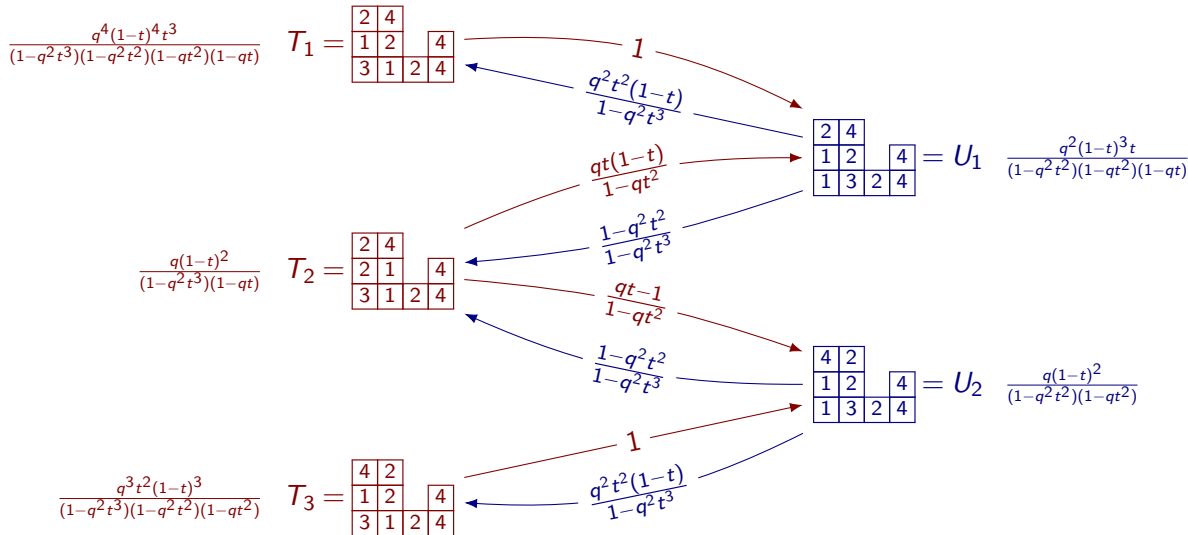
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□

Example (Spoiler)



Part II

Macdonald Polynomials and Non-Attacking Fillings

Symmetric Macdonald polynomials

(Macdonald 1988)

$$P_{\lambda}(x_1, \dots, x_n; q, t)$$

where λ is a partition.

- Unique basis under triangularity and orthogonality conditions.
- Extension of Jack polynomials, Hall–Littlewood polynomials, q -Wittaker polynomials, Schur polynomials.

Nonsymmetric Macdonald polynomials

(Macdonald 1996)

$$E_{\alpha}(x_1, \dots, x_n; q, t)$$

where α is a composition.

- Help understand symmetric Macdonald polynomials.
- Extend Demazure characters (key polynomials) and atom polynomials.

Permuted basement Macdonald polynomials

(Ferreira 2011)

$$E_{\alpha}^{\sigma}(x_1, \dots, x_n; q, t)$$

where α is a composition and σ is a permutation.

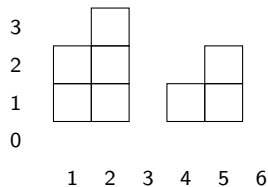
They can be described as:

- a transformation of nonsymmetric Macdonald polynomials by Demazure–Lusztig operators,
- eigenfunctions of a version of Cherednik–Dunkl operators,
- the weight generating functions of **non-attacking fillings**.

Augmented Skyline Diagrams

skyline diagram of a composition α :

$$\alpha = (2 \ 3 \ 0 \ 1 \ 2 \ 0)$$

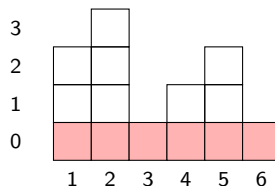


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Augmented skyline diagram of a composition α :

(basement = row 0 = red)

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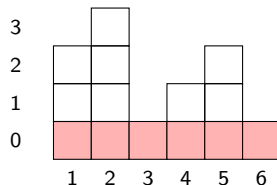


Augmented Skyline Diagrams

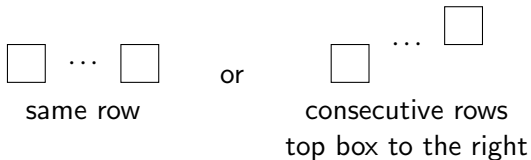
Augmented skyline diagram of a composition α :

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A pair of boxes is **attacking** if:



Non-Attacking Fillings

A **non-attacking filling** of **shape α** is an assignment of numbers $\{1, \dots, n\}$ to the **diagram of α** such that attacking boxes have different entries.

e.g. shape $\alpha = (2, 2, 0, 1)$

2	4		
1	2		4
3	1	2	4

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The **basement** of a filling is the permutation $\sigma = (\sigma_1, \sigma_2, \dots, \sigma_n) \in S_n$ of the entries in the basement row, from left to right.

e.g. shape $\alpha = (2, 2, 0, 1)$ and basement $\sigma = [3, 1, 2, 4]$

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The **content** counts how many i 's are in the filling (excluding the basement).

e.g. the filling above has **content** $(1, 2, 0, 2)$

$$x^T = x_1^1 x_2^2 x_3^0 x_4^2$$

Statistics of Non-Attacking Fillings: maj and coinv

- major index:
- coinversion number:

Statistics of Non-Attacking Fillings: maj and coinv

- major index: $\text{maj}(T) = \sum_{T(u) > T(\text{south of } u)} (1 + \text{leg}(u)),$

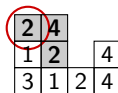
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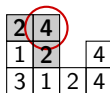


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- coinversion number: $\text{coinv}(T) = \#\{\text{coinversion triples}\}$

2	4		
1	2		4
3	1	2	4

$$1 < 2 < 4$$

2	4		
1	2		4
3	1	2	4

$$1 < 2 < 3$$

2	4		
1	2		4
3	1	2	4

$$4 > 3 > 1$$

2	4		
1	2		4
3	1	2	4

$$1 < 2 < 4$$

2	4		
1	2		4
3	1	2	4

$$4 \geq 4 > 2$$

$$\text{coinv}(T) = 3$$

Weight of a Non-Attacking Filling

$$\text{wt}_{q,t}(T) = q^{\text{maj}(T)} t^{\text{coinv}(T)} \prod_{T(u) \neq T(\text{south of } u)} \frac{1-t}{1 - q^{1+\text{leg}(u)} t^{1+\text{arm}(u)}}$$

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$$\text{wt}_{q,t}(T) = q^4 t^3 \frac{(1-t)^4}{(1-q^2 t^3)(1-q^2 t^2)(1-qt^2)(1-qt)}.$$

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Generating function formula for E_{α}^{σ} (Ferreira 2011)

$$E_{\alpha}^{\sigma} = \sum_{T \in \text{NAF}(\alpha, \sigma)} x^T \text{wt}_{q,t}(T).$$

Alexandersson (2019) studies **symmetries of permuted basement Macdonald polynomials**.

- One of the results involves E_{λ}^{σ} where σ is the **shortest** permutation that sends λ to a rearrangement α .

Corteel, Mandelshtam, and Williams (2022) study the **asymmetric simple exclusion process (ASEP), a model of interacting particles**.

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Theorem (D–Mandelshtam, 2025⁺)

$$E_\lambda^\sigma = E_\lambda^\tau$$

where σ (resp. τ) is the shortest (resp. longest) permutation that sends λ to a rearrangement α .

More generally, we have the following result.

Theorem (D–Mandelshtam, 2025⁺)

Let α be a composition with $\alpha_i = \alpha_{i+1}$ and σ be a permutation. Then,

$$E_{\alpha}^{\sigma} = E_{\alpha}^{\sigma s_i},$$

where $\sigma s_i = [\sigma_1, \dots, \sigma_{i+1}, \sigma_i, \dots, \sigma_n]$.

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How to prove it? **Construct a probabilistic bijection** between $\text{NAF}(\alpha, \sigma)$ and $\text{NAF}(\alpha, \sigma s_i)$ when $\alpha_i = \alpha_{i+1}$.

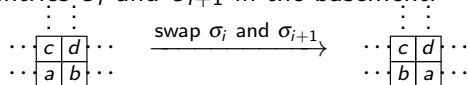
Part III

Probabilistic Bijections for Non-Attacking Fillings

Goal

Construct a probabilistic bijection between $\text{NAF}(\alpha, \sigma)$ and $\text{NAF}(\alpha, \sigma s_i)$, where $\alpha_i = \alpha_{i+1}$.

- **Naive idea:** Just swap the entries σ_i and σ_{i+1} in the basement.



- **An issue:** The resulting filling may be attacking.
- **Idea for fix:** If the resulting filling is attacking, swap entries in the top row as well, and so on.
- **Refinement:** To take care of the weights, assign probabilities to each step.

Definition of the Probabilistic Algorithm by Example

$$\alpha = (3, 4, 4, 0, 0),$$

$$i = 2$$

$$\sigma = [5, 1, 3, 4, 2]$$

$$\sigma s_i = [5, 3, 1, 4, 2]$$

Input:

	3	2		
4	1	2		
3	4	1		
3	4	2		
5	1	3	4	2

Definition of the Probabilistic Operator by Example

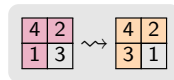
INPUT

	3	2
	1	2
	4	1
	4	2
+	1	3

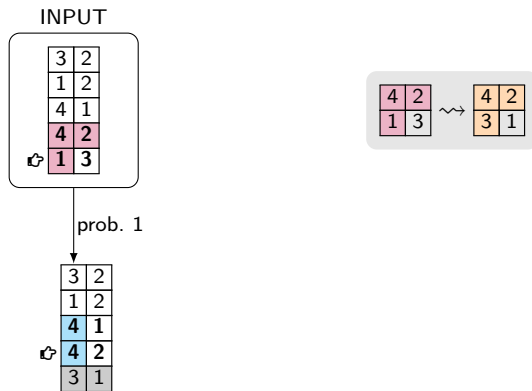
Definition of the Probabilistic Operator by Example

INPUT

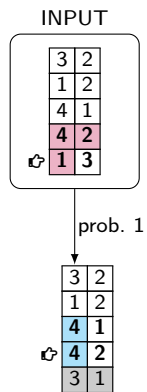
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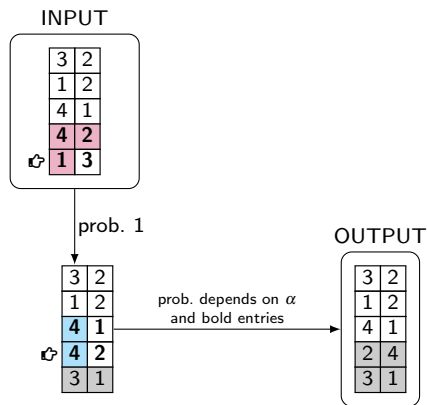
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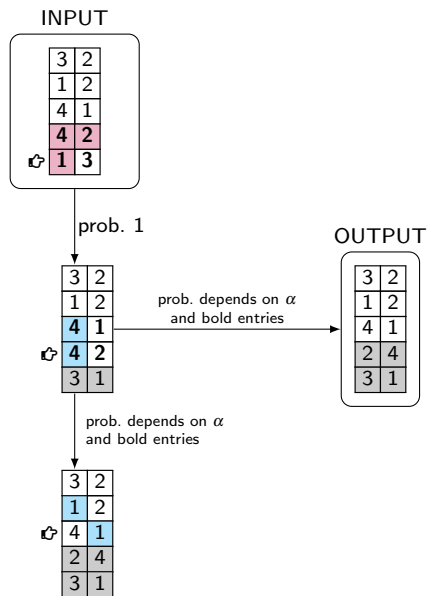
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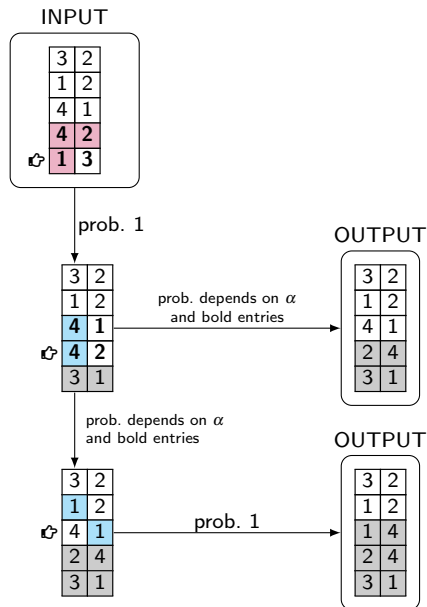
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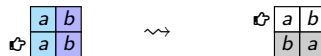
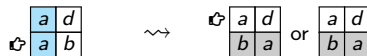
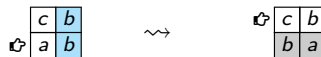
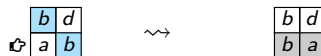
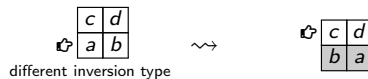
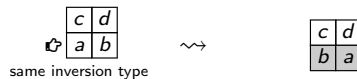
Definition of the Probabilistic Operator by Example



Definition of the Probabilistic Operator by Example



All rules about moving the pointer



It is a probabilistic bijection!

Theorem (D–Mandelshtam, 2025⁺)

The algorithm describes a probabilistic bijection between $\text{NAF}(\alpha, \sigma)$ and $\text{NAF}(\alpha, \sigma s_i)$.

Summary

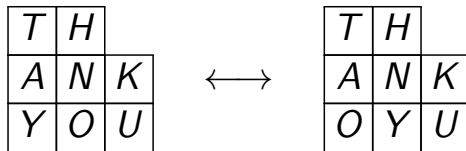
- Probabilistic bijections generalize weight-preserving bijections.
- Write the permuted basement Macdonald polynomial E_{α}^{σ} as the generating function of non-attacking fillings $\text{NAF}(\alpha, \sigma)$.
- Construct a probabilistic bijection between $\text{NAF}(\alpha, \sigma)$ and $\text{NAF}(\alpha, \sigma s_i)$ when $\alpha_i = \alpha_{i+1}$.
- Conclude that, when $\alpha_i = \alpha_{i+1}$,

$$E_{\alpha}^{\sigma} = \sum_{T \in \text{NAF}(\alpha, \sigma)} \text{wt}(T) = \sum_{U \in \text{NAF}(\alpha, \sigma s_i)} \text{wt}(U) = E_{\alpha}^{\sigma s_i}.$$

Summary

- Probabilistic bijections generalize weight-preserving bijections.
- Write the permuted basement Macdonald polynomial E_{α}^{σ} as the generating function of non-attacking fillings $\text{NAF}(\alpha, \sigma)$.
- Construct a probabilistic bijection between $\text{NAF}(\alpha, \sigma)$ and $\text{NAF}(\alpha, \sigma s_i)$ when $\alpha_i = \alpha_{i+1}$.
- Conclude that, when $\alpha_i = \alpha_{i+1}$,

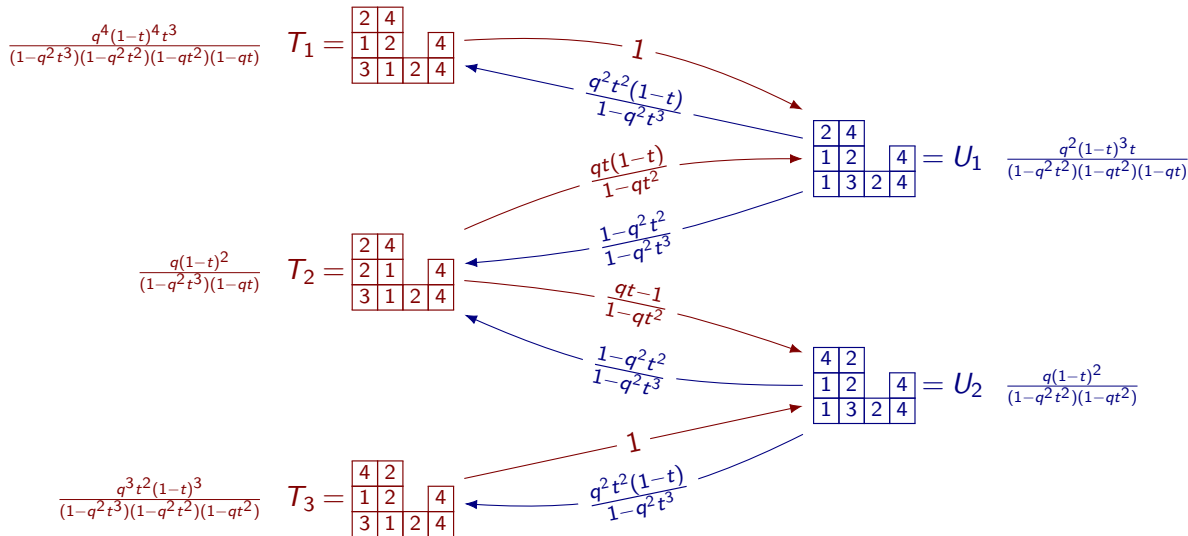
$$E_{\alpha}^{\sigma} = \sum_{T \in \text{NAF}(\alpha, \sigma)} \text{wt}(T) = \sum_{U \in \text{NAF}(\alpha, \sigma s_i)} \text{wt}(U) = E_{\alpha}^{\sigma s_i}.$$



Part IV

Appendix

$\alpha = (2, 2, 0, 1)$, $\sigma = [3, 1, 2, 4]$, content $(1, 2, 0, 2)$



Future Work: $\alpha_i \neq \alpha_{i+1}$ via probabilistic bijections

If $\alpha_i < \alpha_{i+1}$ and $\sigma_i > \sigma_{i+1}$, then

$$E_{s_i \alpha}^{\sigma}(X; q; t) = E_{\alpha}^{\sigma s_i}(X; q; t) + \frac{t^{\phi(\alpha, \sigma, i)}(1-t)}{1-q^{L+1}t^A} E_{\alpha}^{\sigma}(X; q; t)$$

where $A = \text{arm}(x)$ and $L = \text{leg}(x)$ for $x = (i+1, \alpha_i+1) \in \text{dg}(\alpha)$.

Idea for Proof of Balance Condition

- Define the component of the weight contributed by the row r such that

$$\text{wt}_{q,t}(T) = \prod_{\text{row } r} \text{wt}_{q,t}^{(r)}(T).$$

- Define the probability that the pointer moves up in row r denoted by $\rho^{(r)}(T)$.
- Prove “balance condition” row-by-row:

$$\begin{aligned} \text{wt}_{q,t}^{(r)}(T) \cdot \rho^{(r)}(T) &= \text{wt}_{q,t}^{(r)}(U) \cdot \rho^{(r)}(U) \quad \text{or} \\ \text{wt}_{q,t}^{(r)}(T) \cdot (1 - \rho^{(r)}(T)) &= \text{wt}_{q,t}^{(r)}(U) \cdot (1 - \rho^{(r)}(U)), \end{aligned}$$

with the appropriate choice of ρ (move pointer up) or $1 - \rho$ (delete).

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